



Asymmetry in PTP Networks

- What it is and what do we do about it

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What is Asymmetry?

- ❑ **Most packet network links are bidirectional**
 - Traffic can flow from A to B ...
 - ... and from B to A
- ❑ **Asymmetry is simply the difference in transit time between the two directions**
- ❑ **Since a packet only goes in one direction asymmetry cannot be measured on a per-packet basis**
- ❑ **Instead, asymmetry is measured over time**
 - From a few microseconds to a few days or longer
 - Packet times are averaged to determine asymmetry
 - Maybe after some appropriate selection criteria are applied



Why Asymmetry Matters to Sync

❑ It never used to ...

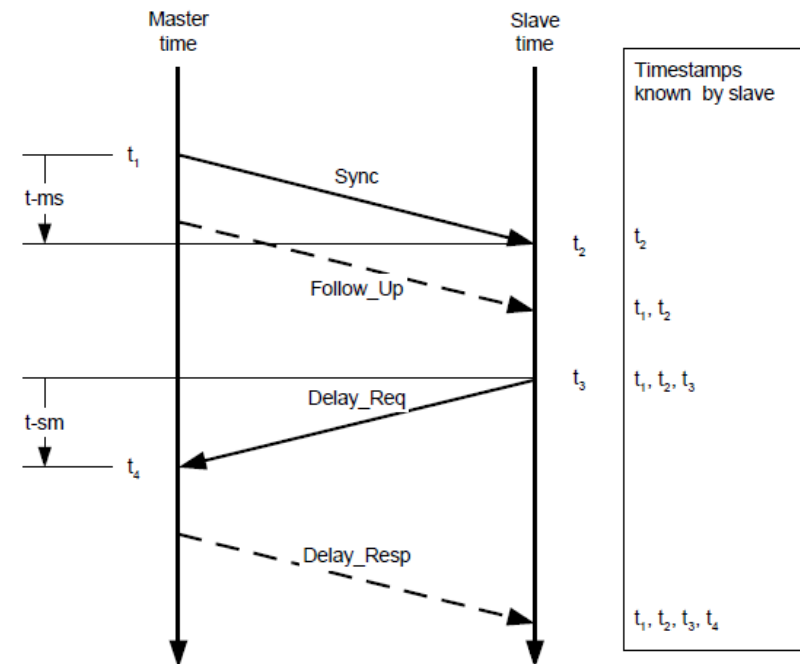
- Back in the days of frequency-only delivery
- Packet-based frequency delivery mechanisms such as adaptive clock recovery can syntonize B to A using just a packet flow from A to B
- Since only one direction is involved asymmetry is irrelevant

❑ However, we now have to deliver time

- Two stage process
 - Syntonize B to A
 - Measure delay from A to B to know how much to adjust B's clock by
- Syntonization is no different to frequency-only case
- Problem is with the delay measurement
 - Chicken and egg situation – to measure delay of packet between A and B requires A and B are synchronized but the whole point of the measurement is to allow synchronization
 - Solution is to measure roundtrip delay of A to B and back to A (or vice-versa)
 - Works because all measurements now at one end
 - However, synchronization requires A-B delay, not A-B-A delay
 - Only solution is to assume A-B delay is half of A-B-A delay
 - Asymmetry breaks this assumption

Synchronization in PTP

- ❑ PTP uses sync messages and delay request messages to measure round-trip delay
 - And delay response message to get t_4 timestamp back to slave
- ❑ Master to slave delay is calculated as:
 - $\text{Delay} = [(t_4 - t_1) - (t_3 - t_2)] / 2$
 - Subtracts time slave took to respond to sync message from roundtrip delay seen by master
 - Gives roundtrip delay of network
 - Divide-by-two gives master to slave delay
 - **Assuming no asymmetry**
 - Any asymmetry results in a phase offset of one half the asymmetry



Defining Asymmetry

- ❑ **Not all asymmetry is made equal**
- ❑ **We can define three types of asymmetry**
 - Based on their duration
 - Or how often they change
- ❑ **Dynamic asymmetry**
 - Changes over a relatively short timeframe
 - As low as few milliseconds
 - To as much as a few minutes
 - Of the same order as the slave's filtering bandwidth
- ❑ **Static asymmetry**
 - Constant – never changes
- ❑ **Pseudo-static asymmetry**
 - The twilight zone between static and dynamic
 - Constant for a lengthy period but not forever

▼ **Bandwidth**



Causes of Static Asymmetry

❑ Static asymmetry is typically caused by network design

❑ Cause 1 : Mismatched fiber lengths

- Separate fibers typically used for each direction
- Often little attention paid to matching fiber lengths
- Or lengths deliberately different due to dispersion compensation
- This will lead to asymmetry
 - Speed of light in fiber is about half of that in a vacuum
 - ~ 6 inches per nanosecond
 - A 200 foot difference in fiber lengths gives 400 ns asymmetry (200 ns phase error)

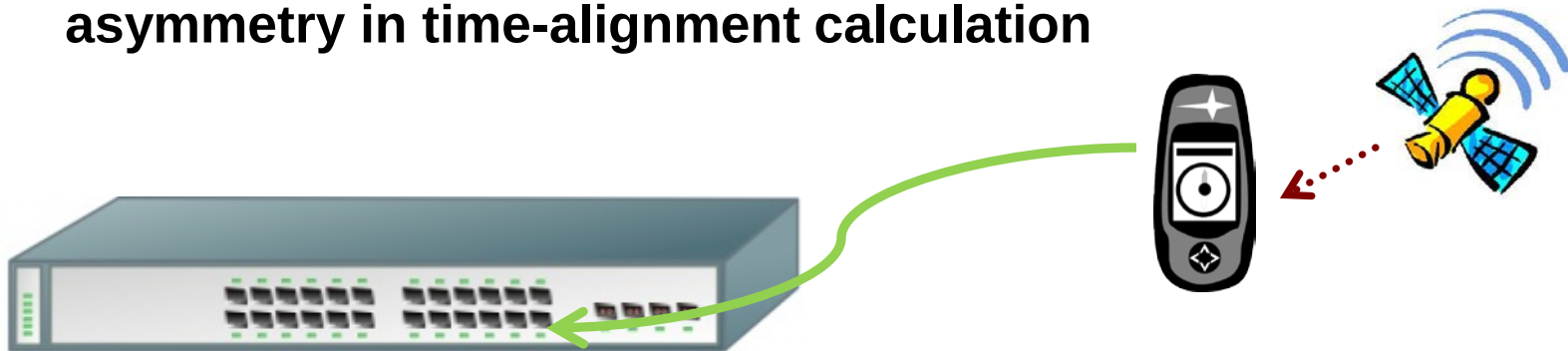
❑ Cause 2: Master and slave speed mismatch

- Occurs when master and slave Ethernet ports are different speed
 - For example 1Gbps master connected to 100Mbps slave
- Results from store-and-forward nature of most Ethernet switches
 - Packet cannot move through switch until entire packet received
 - Cancels out everywhere except end points
- 100Mbps / 1Gbps mismatch gives asymmetry of around 8 μ s

❑ Other smaller static asymmetries possible through poor equipment design

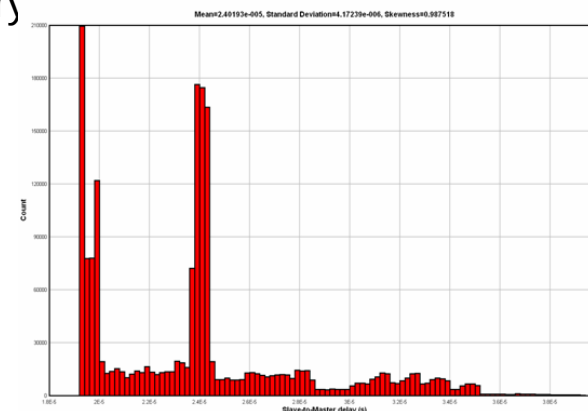
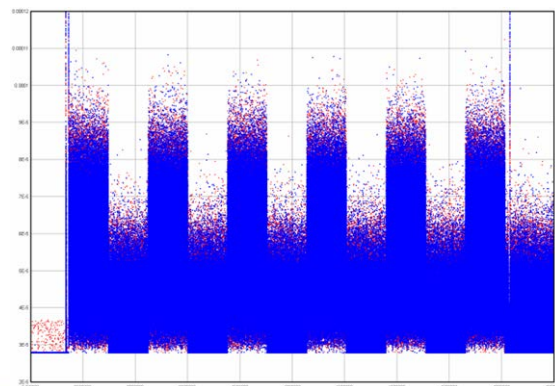
Dealing with Static Asymmetry

- ❑ **The good news ...**
 - Static asymmetry is static
- ❑ **Therefore it can be eliminated forever if the value is known**
- ❑ **Either calculate the asymmetry**
 - For example, resulting from speed mismatch or equipment design
- ❑ **Or measure it**
 - Either independently of the sync solution or with some sort of probe such as a portable GPS receiver
- ❑ **Once measured, program the sync slave to include the known asymmetry in time-alignment calculation**



Causes of Dynamic Asymmetry

- ❑ **Dynamic asymmetry is almost always caused by congestion**
- ❑ **PTP event messages are competing with other network traffic**
- ❑ **Queue blocking occurs within switches and routers**
 - Competition to get both into and out of the egress queue
 - Quality-of-service configuration can partly reduce this by allowing PTP messages to jump ahead of everyone else waiting to egress
 - But QoS does nothing to help getting out of the egress queue
 - Causes variation in packet delay as network traffic changes
- ❑ **Many applications have very asymmetric traffic flows**
 - Think of someone watching YouTube on their iPhone
 - This leads to a traffic-induced asymmetry



Dealing with Dynamic Asymmetry

❑ **Dynamic asymmetry is by definition short-term**

- Therefore, the time recovery slave has a fighting chance of being able to filter it out

❑ **How good a job can be done depends on:**

- Duration and profile of asymmetry
- How smart the slave algorithms are

❑ **Possible algorithmic filtering approaches**

- Packet selection
 - There may be a clear asymmetry when the average transit time of all packets is considered, but slaves don't typically look at every single packet
 - Minimum delay may not actually move at all
- Low-bandwidth filtering
 - If averaging process occurs over much longer period than the asymmetry then the effect will be reduced
 - But puts strain on local oscillator performance
- There are also a number of more advanced techniques
 - For example, it might be possible to detect when the spread of packet delays in one direction has become excessive and temporarily enter a mini-holdover mode

Causes of Pseudo-static Asymmetry

- ❑ **Pseudo-static asymmetry is an asymmetry that could change after an arbitrarily long time**
 - Could change after a network has been operating for weeks or months
 - Or could be a few hours ...
 - ... or a few years
- ❑ **Likely cause is some network architecture change**
 - Instigated by the carrier, or maybe automatic to cover a failure
 - For example, network routing could be changed introducing a new static asymmetry
 - Multiplexing of data onto fiber could be changed
 - ROADM reconfiguration in WDM system
 - Change of wavelength will have effect on packet delay
- ❑ **It may be possible to negotiate a service level agreement that ensures no network changes during normal operation**
 - Could be expensive option though
 - Still doesn't deal with failure conditions

Dealing with Pseudo-static Asymmetry

☐ Can we handle pseudo-static asymmetry in the same way as the others?

- Likely too long-term to filter, so the techniques for dynamic asymmetry won't work
- It's not truly static, so the techniques for static asymmetry won't work
- So, the answer is "no"

☐ The million dollar question:

“What can we do to handle the effects of pseudo-static asymmetry?”

☐ If we can answer this then a major hurdle of PTP deployment has been cleared

☐ If we can't, then PTP will struggle in many applications

- And it will only get worse as timing requirements become tighter

Using Two Timing Sources

❑ The problem is clear

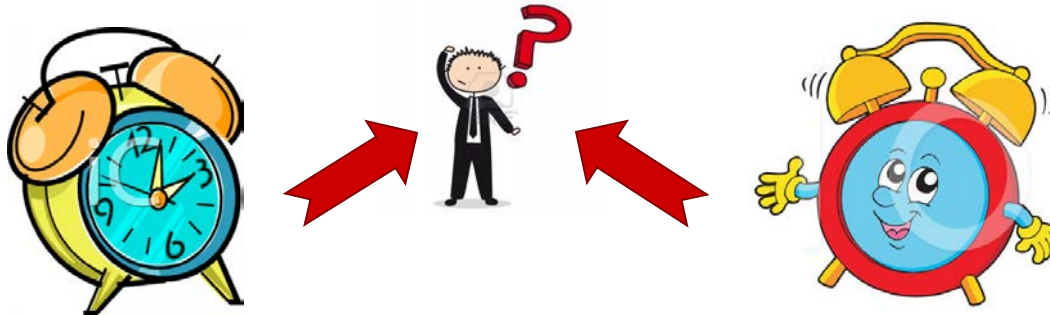
- We have a variable that we cannot predict or measure that impacts the time recovery without the slave knowing what is happening

❑ The only possible solution is to have more than one source of timing

- We can then use one as a reference to judge the performance of the other
- This should allow us to detect both asymmetry changes and other disturbances

❑ Two possible modes:

- Active / active – equal weight is given to both timing sources and the output is derived from some combination of both in normal operation
- Active / probe – the output is derived from a primary timing source with the secondary one just performing a probe function



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Majority Voting

- ☐ **There is actually a third timing source**
 - Time derived from the slave's local oscillator
- ☐ **In an ideal system all three timing sources will remain in sync**
 - No network packet delay variation
 - No network asymmetry
 - Perfect stable local oscillator
- ☐ **In practice, different events will cause the three to diverge**
- ☐ **Consider remote sources A and B and local time L**
 - B and L stay in sync – A jumps with respect to them
 - Likely that network asymmetry to A has changed
 - A and L stay in sync – B jumps with respect to them
 - Likely that network asymmetry to B has changed
 - A and B stay in sync – L moves with respect to them
 - Likely that the local oscillator has drifted
- ☐ **Once we can judge what has changed we can correct for it**
 - For example, if time from A jumps with respect to B and L we can use difference to correct a known asymmetry value and bring A back into line

PTP + Synchronous Ethernet

☐ Synchronous Ethernet is frequency only

- But it's a stable frequency
- Can be used to maintain a local short term time-holdover clock

☐ Maintain two clocks synchronized to PTP

- One higher bandwidth locked directly to PTP
- Second clock lower bandwidth using SyncE frequency as “tick”
 - Effectively a local time holdover clock

☐ During normal operation both clocks will remain the same

- But an asymmetry change or disturbance will cause the two to differ
- If the change appears constant enough over the period in which we have confidence in the SyncE clock then the difference can be used as an asymmetry correction

☐ This is really just a smart use of a PTP / SyncE hybrid mode

- How well it works depends on the system architecture
- Congruent mode – not so good
- Coherent mode – possibility of excellent immunity to network changes

Using PTP + GPS

- ☐ **Everyone knows that GPS is a good source of time ...**
- ☐ **But, GPS availability falls far short of that needed for telecoms**
 - Spoofing
 - Deliberate or accidental jamming
 - Multi-path effects
 - Weather
- ☐ **However, GPS is still good enough to be the primary sync source as long as it has a backup**
- ☐ **PTP can be used as that back-up**
 - As a remote alternative to a highly stable local oscillator
- ☐ **The GPS can be used for a continual asymmetry measurement during normal operation**
- ☐ **PTP takes over in the event of a GPS outage**
 - With the likely asymmetry already known
- ☐ **Easily implemented by replacing existing GPS holdover module with PTP-equipped version**

Using Dual PTP Flows

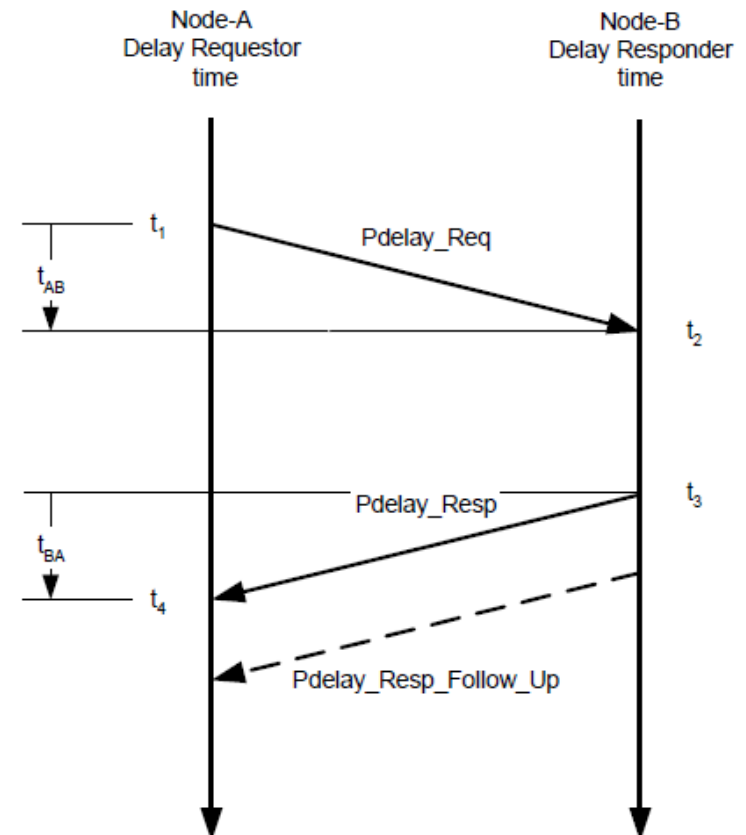
- ❑ **Instead of PTP + GPS, how about PTP + PTP?**
- ❑ **Need to lock to two separate masters simultaneously**
 - Or could lock twice to the same master
- ❑ **Path from slave to each master must be independent**
 - Otherwise asymmetry change on shared portion won't be detected
 - Can relax this for local connection since we're in control
- ❑ **This is similar to the standard dual-master upstream redundancy model**
 - Except now, the standby source is used to detect asymmetry changes associated with the primary source and vice-versa
- ❑ **Low cost implementation from the slave standpoint**
 - Likely doesn't require any additional hardware
- ❑ **However, very hard to architect in practice**
 - Since totally separate routes to masters are required
 - Could require dual links from independent carriers to make it work

Monitoring an Array of Boundary Clocks

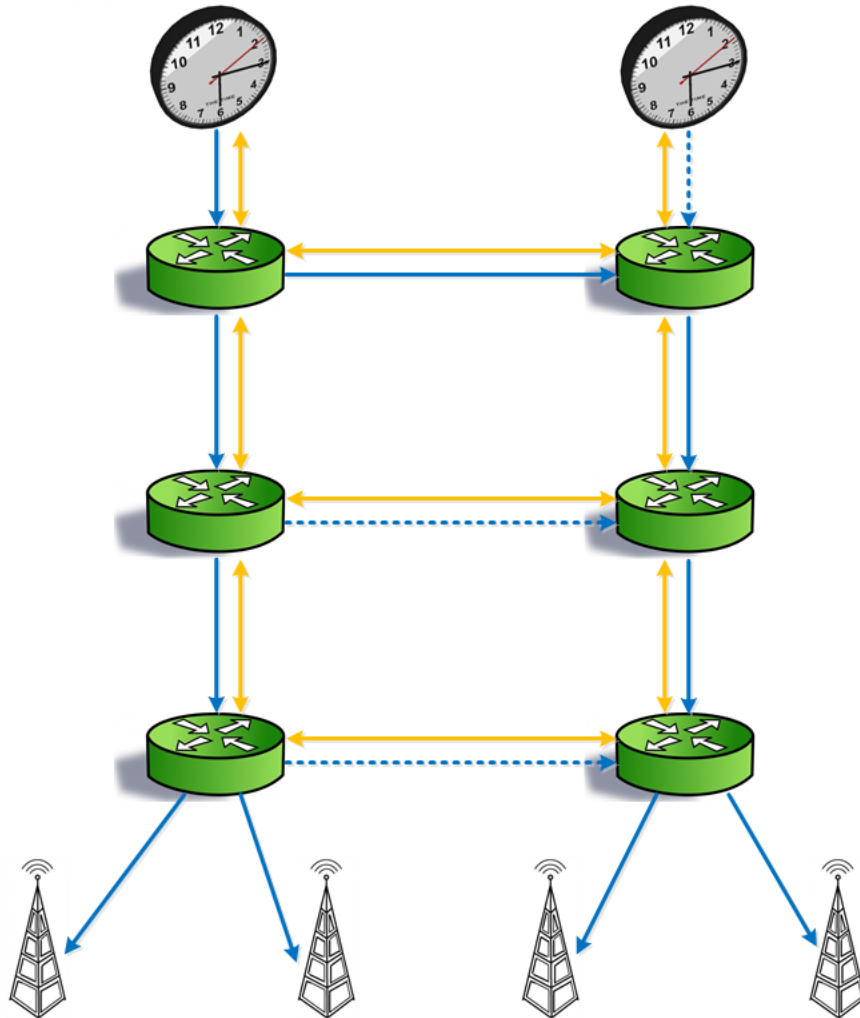
- ☐ It has been proposed that multiple PTP flows could be used to monitor the performance of an array of boundary clocks
 - Asymmetry
 - Congestion effects
 - General operation
- ☐ Assumes that every node in the network is a boundary clock
- ☐ Assumes that multiple masters are available
- ☐ Using the standard IEEE 1588-2008 Best Master Clock Algorithm (BMCA) results in:
 - Only one master being selected by all slaves
 - Each boundary clock receiving its timing from a single boundary clock or grandmaster
- ☐ All superfluous links to boundary clock forced to passive state
 - Not allowed to take part in standard Sync / Delay Request activity
 - But can still run peer-to-peer delay measurement
- ☐ Pdelay mechanism can be used to support secondary timing flow
 - Normally used as part of peer-to-peer transparent clock mechanism

The PTP Pdelay Mechanism

- ❑ **Runs bidirectionally across all links at all times**
 - No master or slave concept
 - Independent of port state
- ❑ **Works through Pdelay_Req and Pdelay_Resp messages**
 - Both timestamped event messages
- ❑ **IEEE 1588-2008 defines various modes involving timestamps and correction fields**
 - Importantly, one mode allows the requestor to get both the responder's time-of-day and the delay
 - Provides a secondary timing source



Example Pdelay Use



- ❑ **Network of multiple masters and slaves**
 - Boundary clocks at each node
- ❑ **Blue lines indicate primary timing flow**
 - Dotted when excluded by BMCA
 - Only one master is active
- ❑ **Orange lines should Pdelay secondary timing flow**
 - Active between each node and its neighbors
 - Allows each node to get a secondary timing source to use as a probe
- ❑ **Remains to be seen how practical this is in a real network**

Summary

- ❑ **Asymmetry impacts packet-based time transport**
 - Will become more of a problem as:
 - PTP deployment becomes more widespread
 - Applications and technologies drive for tighter and tighter sync requirements
- ❑ **We need to worry about three types of asymmetry –**
 - Dynamic – short-term
 - Static – never changes
 - Pseudo-static – changes with a long period
- ❑ **Different techniques needed to handle each type**
 - Dynamic
 - Clever filtering within the slave
 - Smart network architecture – QoS and avoiding congestion
 - Static
 - Measurement
 - Calculation
 - Pseudo-static
 - Multiple timing flows
 - Majority voting
- ❑ **PTP + GPS and PTP + PTP are both possible options**
- ❑ **Use of Pdelay mechanism may be an option too**